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## Droste-Effekt bei präspositiven Zeichenklassen

1. In Toth (2011) hatten wir das System der präspositiven Zeichenrelationen wie folgt dargestellt:

$$\begin{aligned}
 & \left[ (\{\emptyset\}. \{2', 2''\} \{3'\}. \{2', 2''\} \{2', 2''\}. \{2', 2''\}) \times \right. \\
 & \quad \left. (\{2', 2''\}. \{2', 2''\} \{2', 2''\}. \{3'\} \{2', 2''\}. \{\emptyset\}) \right] \\
 & \left[ (\{\emptyset\}. \{2', 2''\} \{3'\}. \{2', 2''\} \{2', 2''\}. \{3'\}) \times \right. \\
 & \quad \left. (\{3'\}. \{2', 2''\} \{2', 2''\}. \{3'\} \{2', 2''\}. \{\emptyset\}) \right] \\
 & \left[ (\{\emptyset\}. \{2', 2''\} \{3'\}. \{2', 2''\} \{2', 2''\}. \{\emptyset\}) \times \right. \\
 & \quad \left. (\{\emptyset\}. \{2', 2''\} \{2', 2''\}. \{3'\} \{2', 2''\}. \{\emptyset\}) \right] \\
 & (\{\emptyset\}. \{2', 2''\} \{3'\}. \{3'\} \{2', 2''\}. \{3'\}) \times (\{3'\}. \{2', 2''\} \{3'\}. \{3'\} \{2', 2''\}. \{\emptyset\}) \\
 & (\{\emptyset\}. \{2', 2''\} \{3'\}. \{3'\} \{2', 2''\}. \{\emptyset\}) \times (\{\emptyset\}. \{2', 2''\} \{3'\}. \{3'\} \{2', 2''\}. \{\emptyset\}) \\
 & (\{\emptyset\}. \{2', 2''\} \{3'\}. \{\emptyset\} \{2', 2''\}. \{\emptyset\}) \times (\{\emptyset\}. \{2', 2''\} \{\emptyset\}. \{3'\} \{2', 2''\}. \{\emptyset\}) \\
 & (\{\emptyset\}. \{3'\} \{3'\}. \{3'\} \{2', 2''\}. \{3'\}) \times (\{3'\}. \{2', 2''\} \{3'\}. \{3'\} \{3'\}. \{\emptyset\}) \\
 & (\{\emptyset\}. \{3'\} \{3'\}. \{3'\} \{2', 2''\}. \{\emptyset\}) \times (\{\emptyset\}. \{2', 2''\} \{3'\}. \{3'\} \{3'\}. \{\emptyset\}) \\
 & (\{\emptyset\}. \{3'\} \{3'\}. \{\emptyset\} \{2', 2''\}. \{\emptyset\}) \times (\{\emptyset\}. \{2', 2''\} \{\emptyset\}. \{3'\} \{3'\}. \{\emptyset\}) \\
 & (\{\emptyset\}. \{\emptyset\} \{3'\}. \{\emptyset\} \{2', 2''\}. \{\emptyset\}) \times (\{\emptyset\}. \{2', 2''\} \{\emptyset\}. \{3'\} \{\emptyset\}. \{\emptyset\})
 \end{aligned}$$

Nun erinnern wir uns, dass gilt

$$CM = \{O', O''\}$$

$$C(M, O) = \{I'\}$$

$$C(M, O, I) = \{\emptyset\},$$

also

$$C(ZR) = C(M, O, I) = (\{O', O''\}, \{I'\}, \{\emptyset\}).$$

Somit erhalten wir wegen

$$\{\emptyset\} = (\{O', O''\}, \{I'\}, \{\emptyset\})$$

in einem 1. Rekursionsschritt

$$\begin{aligned} & \left[ \begin{aligned} & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(2', 2'').(2', 2'').(2', 2'')) \times \\ & (2', 2'').(2', 2'').(2', 2'').(3')\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \end{aligned} \right] \\ & \left[ \begin{aligned} & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(2', 2'').(2', 2'').(3') \times \\ & (3').(2', 2'').(2', 2'').(3')\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \end{aligned} \right] \\ & \left[ \begin{aligned} & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(2', 2'').(2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \times \\ & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \end{aligned} \right] \\ & \left[ \begin{aligned} & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(3')\{2', 2'').(3') \times \\ & \times (3').(2', 2'').(3')\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \end{aligned} \right] \\ & \left[ \begin{aligned} & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(3')\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \times \\ & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(3')\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \end{aligned} \right] \\ & \left[ \begin{aligned} & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \times \\ & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}.(\{O', O''\}, \{I'\}, \{\emptyset\}).\{3'\}\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \end{aligned} \right] \\ & \left[ \begin{aligned} & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{3'\}\{3'\}.(3')\{2', 2'').(3') \times \\ & (3').(2', 2'').(3')\{3'\}\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \end{aligned} \right] \\ & \left[ \begin{aligned} & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{3'\}\{3'\}.(3')\{2', 2'').(\{O', O''\}, \{I'\}, \{\emptyset\})) \times \\ & ((\{O', O''\}, \{I'\}, \{\emptyset\}).\{2', 2''\}\{3'\}\{3'\}.(\{O', O''\}, \{I'\}, \{\emptyset\})) \end{aligned} \right] \end{aligned}$$

$$\begin{aligned}
& \left( (\{\{0', 0''\}, \{I'\}, \{\emptyset\}), \{\{3'\}\{3'\}(\{\{0', 0''\}, \{I'\}, \{\emptyset\})\{2', 2''\}(\{\{0', 0''\}, \{I'\}, \{\emptyset\})) \right) \\
& \times \\
& \left. \left( (\{\{0', 0''\}, \{I'\}, \{\emptyset\}), \{2', 2''\}(\{\{0', 0''\}, \{I'\}, \{\emptyset\}), \{\{3'\}\{3'\}(\{\{0', 0''\}, \{I'\}, \{\emptyset\})) \right) \right. \\
& \left. \left( (\{\{0', 0''\}, \{I'\}, \{\emptyset\}), (\{\{0', 0''\}, \{I'\}, \{\emptyset\})\{3'\}(\{\{0', 0''\}, \{I'\}, \{\emptyset\})\{2', 2''\}(\{\{0', 0''\}, \{I'\}, \{\emptyset\})) \right. \right. \\
& \left. \left. \times \right. \right. \\
& \left. \left. (\{\{0', 0''\}, \{I'\}, \{\emptyset\}), \{2', 2''\}(\{\{0', 0''\}, \{I'\}, \{\emptyset\}), \{\{3'\}(\{\{0', 0''\}, \{I'\}, \{\emptyset\})\{0', 0''\}, \{I'\}, \{\emptyset\}) \right. \right)
\end{aligned}$$

in einem 2. Rekursionsschritt

$$\begin{aligned}
& \left( (\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\})), \{2', 2''\}\{3'\}\{2', 2''\}\{2', 2''\}\{2', 2''\}) \times \right. \\
& \left. (\{2', 2''\}\{2', 2''\}\{2', 2''\}\{3'\}\{2', 2''\}(\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\}))) \right) \\
& \left( (\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\})), \{2', 2''\}\{3'\}\{2', 2''\}\{2', 2''\}\{3'\}) \times \right. \\
& \left. (\{3'\}\{2', 2''\}\{2', 2''\}\{3'\}\{2', 2''\}(\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\}))) \right) \\
& \left( (\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\})), \{2', 2''\}\{3'\}\{2', 2''\}\{2', 2''\}\{2', 2''\}(\{\{0', 0''\}, \{I'\}, \right. \\
& \left. \left. (\{\{0', 0''\}, \{I'\}, \{\emptyset\})) \right) \times \right. \\
& \left. \left. (\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\})), \{2', 2''\}\{2', 2''\}\{3'\}\{2', 2''\}(\{\{0', 0''\}, \{I'\}, \right. \right. \\
& \left. \left. (\{\{0', 0''\}, \{I'\}, \{\emptyset\})) \right) \right. \\
& \left( (\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\})), \{2', 2''\}\{3'\}\{3'\}\{2', 2''\}\{3'\}) \right. \\
& \left. \times (\{3'\}\{2', 2''\}\{3'\}\{3'\}\{2', 2''\}(\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\}))) \right) \\
& \left( (\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\})), \{2', 2''\}\{3'\}\{3'\}\{2', 2''\}\{2', 2''\}(\{\{0', 0''\}, \{I'\}, \right. \\
& \left. \left. (\{\{0', 0''\}, \{I'\}, \{\emptyset\})) \right) \times \right. \\
& \left. \left. (\{\{0', 0''\}, \{I'\}, (\{\{0', 0''\}, \{I'\}, \{\emptyset\})), \{2', 2''\}\{3'\}\{3'\}\{2', 2''\}\{2', 2''\}(\{\{0', 0''\}, \{I'\}, \right. \right. \\
& \left. \left. (\{\{0', 0''\}, \{I'\}, \{\emptyset\})) \right) \right)
\end{aligned}$$

$((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{2', 2''\}.\{3'\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\}))$   
 $\{2', 2''\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\}))) \times$   
 $((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{2', 2''\}.\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{3'\}.\{2', 2''\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})))$   
 $\left. \begin{array}{l} ((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{3'\}.\{3'\}.\{3'\}.\{2', 2''\}.\{3'\}) \\ \{3'\}.\{2', 2''\}.\{3'\}.\{3'\}.\{3'\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\}))) \end{array} \right) \times$   
 $((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{3'\}.\{3'\}.\{3'\}.\{2', 2''\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\}))) \times$   
 $((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{2', 2''\}.\{3'\}.\{3'\}.\{3'\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})))$   
 $\left. \begin{array}{l} ((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{3'\}.\{3'\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\}))) \\ \{2', 2''\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\}))) \end{array} \right) \times$   
 $((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{2', 2''\}.\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{3'\}.\{3'\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})))$   
 $\left. \begin{array}{l} ((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\}))) \\ \{3'\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{2', 2''\}.(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\}))) \end{array} \right) \times$   
 $((\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{2', 2''\}.\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).\{3'\}.\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})).(\{0', 0''\}, \{I'\}, (\{0', 0''\}, \{I'\}, \{\emptyset\})))$

usw.

Man erinnert sich also an ähnliche, durch Zirkularität angelegte und durch Einsetzung systematisch erzeugbare Mirimanoff-Serien beim sog. Droste- oder „La vache qui rit“-Effekt in der Semiotik (Toth 2008). In Übereinstimmung mit den seinerzeit erzielten Ergebnissen halten wir fest: Ganz egal, ob man von einer Zeichendefinition mit Fundierungs- oder Antifundierungs-Axiom ausgeht, das komplementäre System der präsupponierten Zeichenrelationen ist prinzipiell antifundiert.

## **Bibliographie**

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